

Challenge 2

Wednesday, 27 November 2024 01:00

One of the first similarities we can observe between \mathbb{Z} and $F[x]$ is that for any $a(x) \in F[x]$, $a(x) = \beta \Leftrightarrow R_{x-a}(a(x)) = \beta$ ($\deg(x-a)=1 \Rightarrow R_{x-a}(a(x)) \in F$ and $a(x) = q(x)(x-a) + R_{x-a}(a(x)) \Rightarrow a(x) = R_{x-a}(a(x))$)

(This explains why interpolation breaks down when $\deg(a(x)) \geq |F|$)
This brings Lagrange interpolation closer to the Chinese Remainder Theorem, which in the script is defined upon modular congruences, but, as its name may indicate, can easily be reformulated in terms of remainders:

$$\begin{cases} x \equiv_{m_1} a_1 \\ \vdots \\ x \equiv_{m_n} a_n \end{cases} \Leftrightarrow \begin{cases} R_{m_1}(x) = a'_1 \\ \vdots \\ R_{m_n}(x) = a'_n \end{cases} \quad \text{where } a'_i = R_{m_i}(a_i)$$

The definition of this system requires R to be a Euclidean domain, as $R_m(x) = a \Leftrightarrow x = q \cdot m + a$, with a "smaller than" m or $a = 0$.

This can be done with a degree function $d: R \rightarrow \mathbb{N}$.

- (i) $\forall a, b \in R, \exists q, r \in R: a = b \cdot q + r$ with $d(r) < d(b)$
- (ii) $\forall a, b \in R: d(a) \leq d(a, b)$

Furthermore, for the Chinese Remainder Theorem to work, m_1, \dots, m_n need to be coprime. To formalise this notion, we have to introduce a gcd function $\gcd: R^2 \rightarrow R$.

Consider, for any $a, b \in R$, the ideal (a, b) .

$$(a, b) = \{u \cdot a + v \cdot b \mid u, v \in R\}$$

$$a = 1 \cdot a + 0 \cdot b \Rightarrow a \in (a, b) \Rightarrow (a, b) \neq \emptyset$$

Since \mathbb{N} is well-ordered, $\{d(x) \mid x \in (a, b)\}$ has a least element, m .

Define $\gcd(a, b)$ as one of the elements with minimal degree m :

$$d(\gcd(a, b)) = \min_{x \in (a, b)} d(x)$$

We can now define:

$$m \text{ and } n \text{ coprime} \stackrel{\text{def}}{\Leftrightarrow} \gcd(m, n) \sim 1$$

$$\Rightarrow \gcd(m, n) = s \quad \text{with } s \in R^* \text{ (def of } \sim)$$

$$\Rightarrow \text{there exists } u, v \in R \text{ s.t. } u \cdot m + v \cdot n = s (\gcd(m, n) \in (m, n))$$

$$\Rightarrow s^{-1} \cdot u \cdot m + s^{-1} \cdot v \cdot n = 1 \quad (s \in R^*)$$

$$\Rightarrow \exists u', v' \in R: u' \cdot m + v' \cdot n = 1 \quad (\Delta)$$

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uniquely del. ... i takes on a value ... and ... is less than that of.

We now see that n coprime integers, as well as n coprime polynomials of degree 1, uniquely define an integer or a polynomial whose degree is less than that of their product (for integers, we define degree as the absolute value). To be more precise, there exists one equivalence class $E \in R / \prod_{i=1}^n m_i R$ s.t. any $x \in E$ is a solution of the system.

The quotient group R/mR is equal to R/\equiv_m , but proving this isn't necessary, we can just accept R/mR as an alternative notation of R/\equiv_m .

We also define $a \equiv_m b \stackrel{\text{def}}{\iff} m \mid a-b \iff \exists k \in R: a-b = k \cdot m$ (k is denoted $\frac{a-b}{m}$)

For all $a, b, k \in R$:

$$a \mid b+ka \iff a \cdot \frac{(b+ka)}{a} = b+ka \iff a \left(\frac{b+ka}{a} - k \right) = b \implies a \mid b \iff b+ka \equiv_m b \quad (\textcircled{3})$$

Consider for some coprime $m, n \in R$ the function:

$$\varphi: R/mR \times R/nR \rightarrow R/mnR$$

$$\varphi(a, b) = a \cap b$$

$$\gcd(m, n) \sim 1 \implies \text{there exist } u, v \in R \text{ s.t. } um + vn = 1 \quad (\Delta)$$

$$\implies um \equiv_n 1 \wedge vn \equiv_m 1 \quad (\textcircled{2})$$

\implies for any $([x]_m, [y]_n) \in R/mR \times R/nR$:

$$y um + x vn \equiv_m oc \wedge y um + x vn \equiv_n y \quad (\textcircled{3})$$

$$\implies y um + x vn \in [x]_m \cap [y]_n \text{ (def of equivalence class)}$$

$$\implies \varphi([x]_m, [y]_n) \neq \emptyset \quad (\textcircled{4})$$

For any $u \in \varphi([x]_m, [y]_n), v \in R$:

If $v \in [u]_m$:

$$v \equiv_m u \implies v \equiv_m u \wedge v \equiv_m u (a \mid c \implies a \mid c) \implies v \in \varphi([x]_m, [y]_n) \text{ (def of } \varphi)$$

If $v \in \varphi([x]_m, [y]_n)$:

$$v \equiv_m u \wedge v \equiv_n u \implies v \equiv_{mn} u \quad (*) \implies v \in [u]_{mn}$$

$$(*) \gcd(m, n) \sim 1 \implies \text{there exists } k, l \in R \text{ s.t. } km + ln = 1 \quad (\Delta)$$

$$\implies u - v = u - v \cdot (km + ln)$$

$$= km(u-v) + ln(u-v)$$

$$= kmn \frac{u-v}{n} + lmn \frac{u-v}{m} \quad (m \mid u-v \wedge n \mid u-v)$$

$$= mn \left(k \frac{u-v}{n} + l \frac{u-v}{m} \right)$$

$$\implies mn \mid u-v \iff u \equiv_{mn} v$$

$$\implies \varphi([x]_m, [y]_n) = [u]_{mn} \text{ for any } u \in \varphi([x]_m, [y]_n) \quad (\#)$$

$$\implies \varphi([x]_m, [y]_n) \in R/mnR$$

$\implies \varphi$ maps to the codomain

φ is obviously totally and well-defined: the intersection of 2 sets is always defined and unique.

Furthermore, φ is surjective:

For all $[x]_{mn} \in R/mnR$:

$$x \equiv_m x \wedge x \equiv_n x$$

$$\implies x \in [x]_m \cap [x]_n$$

$$\begin{aligned}
 & x \equiv_m x \wedge x \equiv_n x \\
 \Rightarrow & x \in [x]_m \cap [x]_n \\
 \Rightarrow & \varphi([x]_m, [x]_n) = [x]_{mn} \quad (*) \\
 \Rightarrow & \varphi \text{ surjective } ([x]_m, [x]_n) \in R/mR \times R/nR
 \end{aligned}$$

Finally:

For all $([x]_m, [y]_n), ([w]_m, [z]_n) \in R/mR \times R/nR$:

If $\varphi([x]_m, [y]_n) = \varphi([w]_m, [z]_n)$:

There exists some $u \in \varphi([x]_m, [y]_n) \quad (*)$:

$$\begin{aligned}
 & u \equiv_m x \wedge u \equiv_m w \wedge u \equiv_n y \wedge u \equiv_n z \\
 \Rightarrow & x \equiv_m w \wedge y \equiv_n z \quad (\text{transitivity of } \equiv_x) \\
 \Rightarrow & [x]_m = [w]_m \wedge [y]_n = [z]_n \\
 \Rightarrow & ([x]_m, [y]_n) = ([w]_m, [z]_n) \\
 \Rightarrow & \varphi \text{ is injective}
 \end{aligned}$$

φ is thus bijective

\Rightarrow the system $x \equiv_m a \wedge x \equiv_n b$ has exactly one set of solutions in R/mnR . (+)

We now only need to prove the uniqueness of the solution for an arbitrary amount of congruences. We do this by induction.

For all $n \in \mathbb{N}^*$, define the property:

$P(n) \Leftrightarrow \forall m_1, \dots, m_n, a_1, \dots, a_n \in R$:

$\gcd(m_i, m_j) \sim 1$ for all $\{i, j\} \subseteq \{1, \dots, n\}$

$\Rightarrow \exists ! E \in R/\prod_{i=1}^n m_i R : x \in E \Leftrightarrow (\forall i \in \{1, \dots, n\} : x \equiv_{m_i} a_i)$

Trivially, $P(1)$ is true.

Assume $P(k)$ true for some $k \in \mathbb{N}$:

$\forall m_1, \dots, m_{k+1}, a_1, \dots, a_{k+1} \in R$ s.t. $\forall \{i, j\} \subseteq \{1, \dots, k+1\} : \gcd(m_i, m_j) \sim 1$:

$x \equiv_{m_k} a_k \wedge x \equiv_{m_{k+1}} a_{k+1} \Leftrightarrow x \equiv_{m_k m_{k+1}} b$ for some $[b]_{m_k m_{k+1}} \in R/m_k m_{k+1} R$ (+)

$\forall i \in \{1, \dots, k-1\}$:

$\gcd(m_i, m_k) \sim 1 \wedge \gcd(m_i, m_{k+1}) \sim 1$

\Rightarrow There exists $u, v, u', v' \in R$ s.t. $u m_i + v m_k = 1, u' m_i + v' m_{k+1} = 1$ (Δ)

$\Rightarrow (u u' m_i + u v' m_{k+1} + u' v m_k) m_i + v v' m_k m_{k+1} = 1$

$\Rightarrow d(\gcd(m_i, m_k m_{k+1})) \leq d(1) \quad (1 \in (m_k, m_{k+1}))$

$\Rightarrow \gcd(m_i, m_k m_{k+1}) \sim 1$ ($d(1) \leq d(1, 2) + \text{antisymmetry of } \leq$)

By the induction hypothesis, the system is solvable for $m_1, \dots, m_{k-1}, m_k m_{k+1}, a_1, \dots, a_{k-1}, b \Rightarrow P(k+1)$ holds

$m_1, \dots, m_{k-1}, m_k m_{k+1}, a_1, \dots, a_{k-1}, b \Rightarrow P(k+1) \text{ holds}$

We can thus generalise the Chinese Remainder Theorem to any Euclidean domain D :

For any Euclidean domain D , the system of equations S :

$$S: \begin{cases} x \equiv_{m_1} a_1 \\ \vdots \\ x \equiv_{m_n} a_n \end{cases}$$

For some $n \in \mathbb{N}^*$, with $m_1, \dots, m_n, a_1, \dots, a_n \in D$ and $\gcd(m_i, m_j) = 1$ for any $\{i, j\} \subseteq \{1, \dots, n\}$ has exactly one solution $E \in D / \prod_{i=1}^n m_i D$ s.t. $x \in E \Leftrightarrow x$ solves S .